Factor complexity of infinite words associated with β-expansions

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Rényi expansion of unity in base $\beta > 1$

$$d_\beta(1) = t_1 t_2 t_3 \cdots, \qquad t_i = \lfloor \beta T_\beta^{i-1}(1) \rfloor,$$

where

$$T_\beta : [0,1] \to [0,1), \quad T_\beta(x) := \beta x - \lfloor \beta x \rfloor = \{\beta x\}.$$

- Parry number: $d_\beta(1)$ is eventually periodic,
- simple Parry number: $d_\beta(1) = t_1 \cdots t_m$,
- non-simple Parry number: $d_\beta(1) = t_1 \cdots t_m (t_{m+1} t_{m+2} \ldots t_{m+p})^\omega$.

Simple Parry numbers

$d_\beta(1) = t_1 \cdots t_m$

Canonical substitution φ_β over the alphabet $\mathcal{A} = \{0, 1, \ldots, m-1\}$

$$\begin{aligned}
\varphi_\beta(0) &= 0^{t_1}1 \\
\varphi_\beta(1) &= 0^{t_2}2 \\
&\vdots \\
\varphi_\beta(m-2) &= 0^{t_{m-1}}(m-1) \\
\varphi_\beta(m-1) &= 0^{t_m}
\end{aligned}$$

Non-simple Parry numbers

$d_\beta(1) = t_1 \cdots t_m (t_{m+1} t_{m+2} \ldots t_{m+p})^\omega$

Canonical substitution φ_β over the alphabet $\mathcal{A} = \{0, 1, \ldots, m+p-1\}$

$$\begin{array}{rcl} \varphi_\beta(0) & = & 0^{t_1}1 \\ \varphi_\beta(1) & = & 0^{t_2}2 \\ & \vdots & \\ \varphi_\beta(m-1) & = & 0^{t_m}m \\ \varphi_\beta(m) & = & 0^{t_{m+1}}(m+1) \\ & \vdots & \\ \varphi_\beta(m+p-2) & = & 0^{t_{m+p-1}}(m+p-1) \\ \varphi_\beta(m+p-1) & = & 0^{t_{m+p}}m \end{array}$$

Fixed point $\mathbf{u}_\beta = \lim_{n\to\infty} \varphi_\beta^n(0) = 0^{t_1}1\cdots$.

Basic definitions – factor complexity

$\mathcal{A} = \{0, 1, \dots, q-1\}$	alphabet
$\mathbf{u} = (\mathbf{u}_i)_{i\in\mathbb{N}}$, $\mathbf{u}_i \in \mathcal{A}$	infinite word over $\mathcal{A}$
$w = \mathbf{u}_j\mathbf{u}_{j+1}\cdots\mathbf{u}_{j+n-1}$	factor of $\mathbf{u}$ of length n
$\mathcal{L}_n(\mathbf{u})$	the set of factors of $\mathbf{u}$ of length n
$\mathcal{L}(\mathbf{u}) = \bigcup_{n\in\mathbb{N}} \mathcal{L}_n(\mathbf{u})$	the language of $\mathbf{u}$

Factor complexity of $\mathbf{u}$ is the function $\mathcal{C} : \mathbb{N} \to \mathbb{N}$, given by $\mathcal{C}(n) := \#\mathcal{L}_n(\mathbf{u})$.

Basic definitions – fixed point of substitution

$\varphi(0) = 0v$, $v \in \mathcal{A}^+$, then *the fixed point* of φ given by $\mathbf{u} := \lim_{n\to\infty} \varphi^n(0) = \varphi^\omega(0)$ is an infinite word which is *uniformly recurrent*.

A substitution φ is primitive if for all $a, b \in \mathcal{A}$ there exists $k \in \mathbb{N}$ such that the word $\varphi^k(a)$ contains b. In what follows, we assume that φ is *primitive and injective*.

In general, complexity of a fixed point of any primitive substitution is a sublinear function $\mathcal{C}(n) \le an + b$, $a, b \in \mathbb{N}$.

Known results for simple Parry numbers

Simple Parry numbers (Bernat, Frougny, Masáková, Pelantová):

- $t_1 = t_2 = \cdots = t_{m-1}$ or $t_1 > \max\{t_2, \ldots, t_{m-1}\}$ exact value of $\mathcal{C}(n)$ is known,
- in particular, $(m-1)n + 1 \leq \mathcal{C}(n) \leq mn$, for all $n \geq 1$,
- $\mathcal{C}(n)$ is affine $\Leftrightarrow$
  1) $t_m = 1$
  2) for all $i = 2, 3, \ldots, m-1$ we have

$$t_i t_{i+1} \ldots t_{m-1} t_1 \ldots t_{i-1} \preceq t_1 t_2 \ldots t_{m-1}.$$

Then $\mathcal{C}(n) = (m-1)n + 1$.

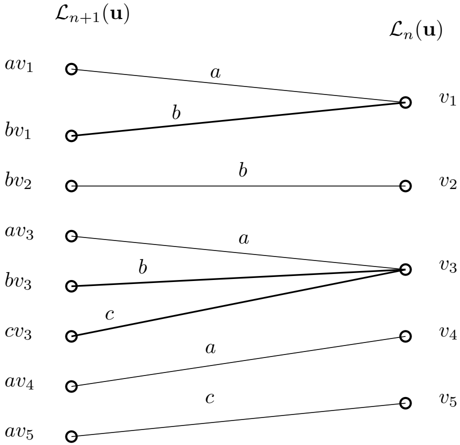
Special factors

For $v \in \mathcal{L}(\mathbf{u})$ we define the set of *left extensions*

$$\mathrm{Lext}(v) := \{a \in \mathcal{A} \mid av \in \mathcal{L}(\mathbf{u})\}.$$

If $\#\mathrm{Lext}(v) > 1$, then v is said to be *left special (LS) factor*.
Analogously are defined right special (RS) factors.

LS factors and factor complexity



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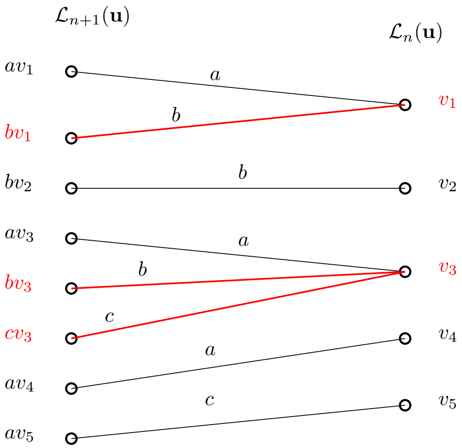

$$\mathcal{C}(n+1) - \mathcal{C}(n) = 3$$

LS factors and factor complexity

For the first difference of the complexity function holds:

$$\triangle\mathcal{C}(n) := \mathcal{C}(n+1) - \mathcal{C}(n) = \sum_{\substack{v\in\mathcal{L}_n(\mathbf{u}) \\ v \text{ is LS}}} (\#\mathrm{Lext}(v) - 1).$$

Complete knowledge of all LS factors along with the number of their left extensions allow us to evaluate $\mathcal{C}(n)$.

$$\triangle\mathcal{C}(n) \geq 1 \text{ for all } n \in \mathbb{N} \Leftrightarrow \mathbf{u} \text{ is aperiodic.}$$

Structure of LS factors – infinite LS branches

Definition

An infinite word $\mathbf{w}$ *is called infinite LS branch of* $\mathbf{u}$ *if each prefix of* $\mathbf{w}$ *is a LS factor of* $\mathbf{u}$.

$$Lext(\mathbf{w}) = \bigcap_{v \text{ prefix } \mathbf{w}} Lext(v).$$

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- **u** periodic $\Rightarrow$ no infinite LS branches,
- **u** aperiodic $\Rightarrow$ at least one infinite LS branch,
- **u** is a fixed point of a primitive substitution $\Rightarrow$ finite number of infinite LS branches
(consequence of the fact that $\triangle\mathcal{C}(n)$ is bounded (Mossé, Cassaigne))

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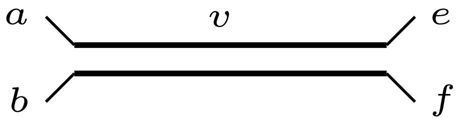

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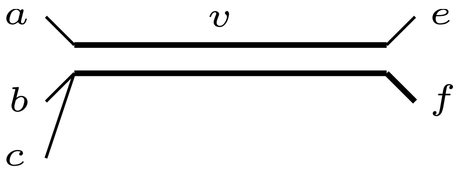

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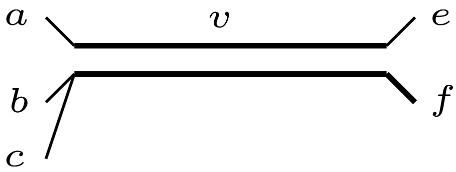

Definition

A LS factor v having $a, b \in Lext(v)$ is called (a, b)-maximal LS factor if for each letter $e \in \mathcal{A}$ we is not a LS factor with left extensions a and b.

Images of LS factors

Example: $\varphi : 1 \mapsto 1211, 2 \mapsto 311, 3 \mapsto 2412, 4 \mapsto 435, 5 \mapsto 534$

$\mathbf{u} = \varphi^{\omega}(1)$

w is a LS factor of $\mathbf{u}$ with left extensions 1 and 2

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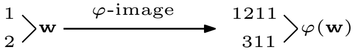

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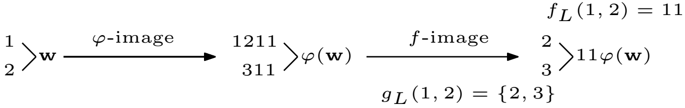

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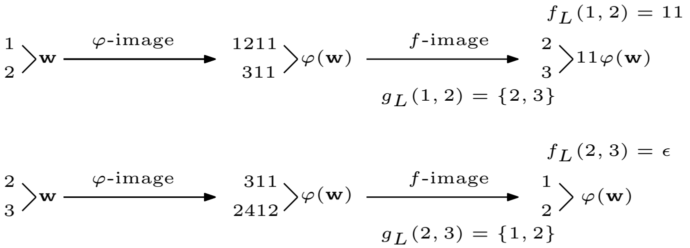

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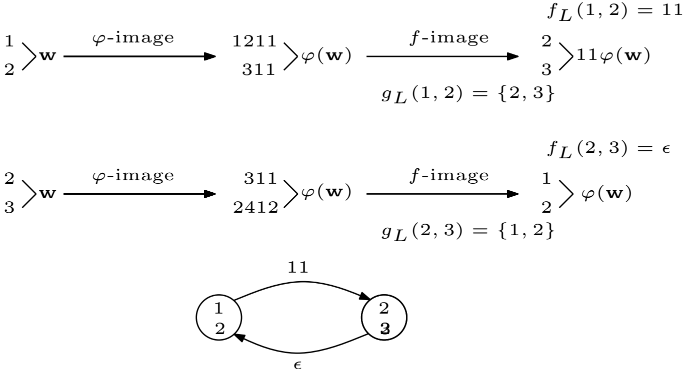

Graph GL_φ

Vertices: unordered couples of distinct letters (a, b).

Edges: if $g_L(a, b) = \{c, d\}$, then there is an edge between (a, b) and (c, d) with label $f_L(a, b)$.

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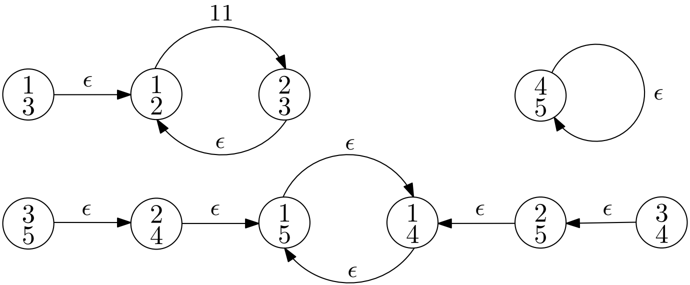
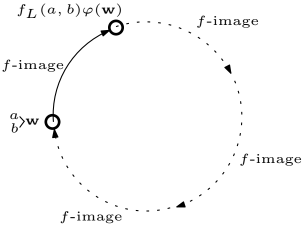


Structure of infinite branches

Assumption: For each infinite LS branch **w** it holds that

a) *f*-image of **w** is uniquely given,

b) there exists exactly one infinite LS branch **w**′ such that **w** is *f*-image of **w**′.



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Theorem

Let $\mathbf{w}$ be an infinite LS branch, $a, b \in Lext(\mathbf{w})$. Then there exists $l > 0$ such that

$$\mathbf{w} = f_L(g_L^{l-1}(a,b)) \cdots \varphi^{l-2}(f_L(g_L(a,b))\varphi^{l-1}(f_L(a,b))\varphi^{l}(\mathbf{w}).$$

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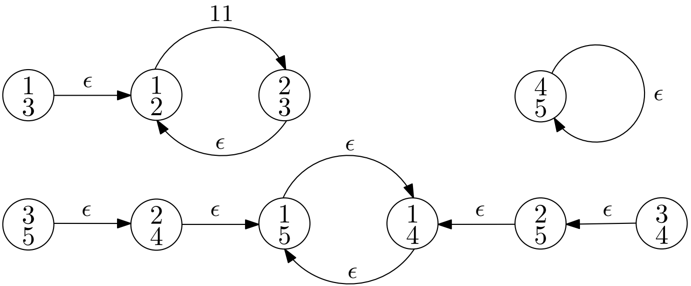
- $f_L = \epsilon \Rightarrow \mathbf{w} = \varphi^l(\mathbf{w})$ and (a,b) is a vertex of a cycle labelled by ϵ only,
- otherwise, (a,b) is a vertex of a cycle labelled not only by ϵ.

Example – how to identify infinite LS branche

$\varphi : 1 \mapsto 1211, 2 \mapsto 311, 3 \mapsto 2412, 4 \mapsto 435, 5 \mapsto 534$

$\mathrm{Lext}(1) = \{1, 2, 3, 4, 5\}, \mathrm{Lext}(2) = \{1, 4, 5\}, \mathrm{Lext}(3) = \{1, 4, 5\}, \mathrm{Lext}(4) = \{1, 2, 3\}, \mathrm{Lext}(5) = \{1, 2, 3\}$



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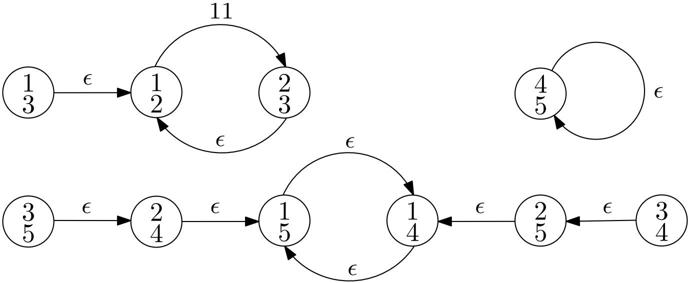

- $\mathbf{w} = 11\varphi^2(\mathbf{w}) \quad \rightarrow \quad 11\varphi^2(11)\varphi^4(11)\cdots$

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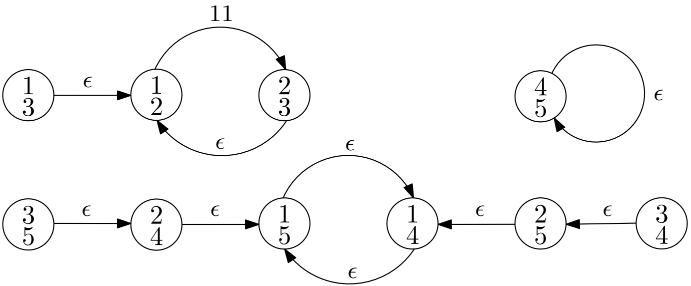

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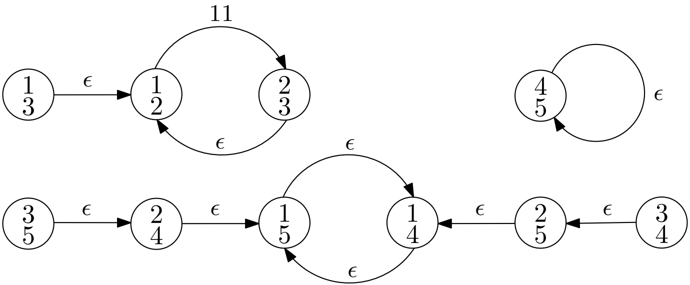

- $\varphi(11)\varphi^3(11)\cdots, \quad 11\varphi^2(11)\varphi^4(11)\cdots$
- $\varphi^\omega(1), \varphi^\omega(4), \varphi^\omega(5), (\varphi^2)^\omega(2), (\varphi^2)^\omega(3)$

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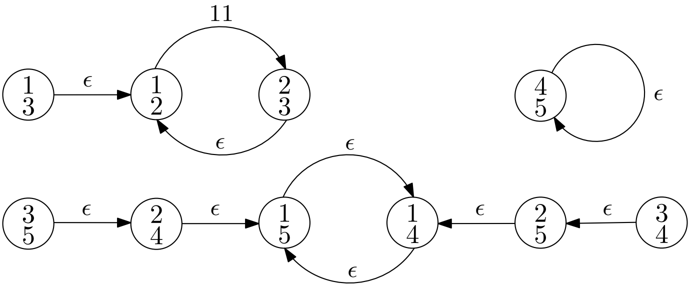

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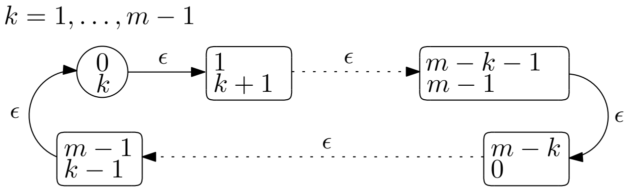
GL_{φ_β} for simple Parry numbers

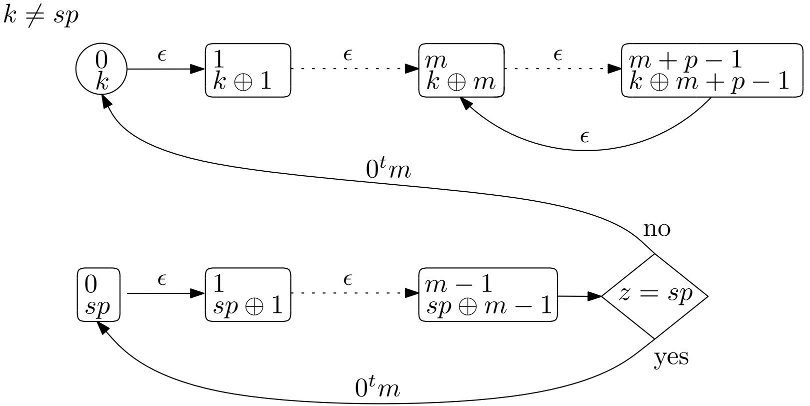
$f_L(a,b) = \epsilon$ for all $a, b \in \{0, 1, \ldots, m-1\}$ and $\mathbf{u}_\beta = \varphi_\beta^\omega(0)$ is the only fixed point



$\Rightarrow$ $\mathbf{u}_\beta$ is the only infinite LS branch

GL_{φ_β} for non-simple Parry numbers

$m-1 \mapsto 0^{t_m} m$, $m+p-1 \mapsto 0^{t_{m+p}} m$, $f_L(m-1, m+p-1) = 0^t m$, $t = \min\{t_m, t_{m+p}\}$, $\mathrm{Lext}(0^t m) = \{0, z\}$, $s \geq 1$



Infinite LS factors

$t = \min\{t_m, t_{m+p}\}$, $\mathrm{Lext}(0^t m) = \{0, z\}$, $s \geq 1$

Definition

$\beta \in \mathcal{S} \Leftrightarrow z = sp \Leftrightarrow$

a) $d_\beta(1) = t_1 \ldots t_m(0 \ldots 0 t_{m+p})^\omega$ and $t_m > t_{m+p}$

b) $d_\beta(1) = t_1 \ldots \underbrace{t_{m-qp}}_{\neq 0} \underbrace{0 \ldots 0}_{qp-1} t_m(t_m + 1 \cdots t_{m+p})^\omega$, $q \geq 1$, $t_m < t_{m+p}$,

$\beta \in \mathcal{S}_0 \Leftrightarrow d_\beta(1) = t_1(0 \cdots 0(t_1 - 1))^\omega$.

Infinite LS factors

Theorem

- *If* β *is a non-simple Parry and* $p > 1$*, then* $\mathbf{u}_\beta$ *is an infinite LS branch with left extensions* $\{m, m+1, \ldots, m+p-1\}$*.*

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Theorem

- *If β is a non-simple Parry and $p > 1$, then $\mathbf{u}_\beta$ is an infinite LS branch with left extensions $\{m, m+1, \ldots, m+p-1\}$.*
- *If $\beta \notin \mathcal{S}$, then $\mathbf{u}_\beta$ is the only one infinite LS branch.*

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- *If β is a non-simple Parry and $p > 1$, then $\mathbf{u}_\beta$ is an infinite LS branch with left extensions $\{m, m+1, \dots, m+p-1\}$.*
- *If $\beta \notin \mathcal{S}$, then $\mathbf{u}_\beta$ is the only one infinite LS branch.*
- *If $\beta \in \mathcal{S}$, then there are m infinite LS branches*

$$0^t m \varphi^m(0^t m)\varphi^{2m}(0^t m)\dots$$
$$\vdots$$
$$\varphi^{m-1}(0^t m)\varphi^{2m-1}(0^t m)\varphi^{3m-1}(0^t m)\dots.$$

Maximal LS factors

f-image of maximal factors

$$\begin{matrix} a \\ b \end{matrix} \rangle \mathbf{w} \langle \begin{matrix} c \\ d \end{matrix} \xrightarrow{f\text{-image}} g_L(a,b) \rangle f_L(a,b)\varphi(\mathbf{w})f_R(c,d) \langle g_R(c,d)$$

Theorem

- *If* $t_1 > 1$ *and* $\beta \notin \mathcal{S}_0$*, then* (a,b)*-maximal factors are* f*-images of the* $(0,p)$*-maximal factor* 0^{t_1-1}

$$\varphi^n(0^{t_1-1}1)(1+n)^{-1}, \quad n = 0,1,\ldots,m-1$$
$$0^t m\varphi^m(0^{t_1-1}1)(1+m)^{-1}, \quad \cdots$$

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$$\begin{aligned} &\varphi^n(0^{t_1-1}1)(1+n)^{-1}, \quad n = 0,1,\ldots,m-1 \\ &0^t m \varphi^m(0^{t_1-1}1)(1+m)^{-1}, \quad \cdots \end{aligned}$$

- *If* $\beta \in \mathcal{S}_0$*, there are no* (a,b)*-maximal factors in* $\mathbf{u}_\beta$*.*

Affine complexity

Theorem

- *The factor complexity of* $\mathbf{u}_\beta$ *is affine* $\Leftrightarrow$ $\mathbf{u}_\beta$ *does not contain any* (a,b)*-maximal factor* $\Leftrightarrow \beta \in \mathcal{S}_0 \Leftrightarrow d_\beta(1) = t_1(0\cdots 0(t_1-1))^\omega$*.* *Then* $\mathcal{C}(n) = (m+p-1)n+1$*.*
  - *The first equivalence is not valid in general (Chacon),*
  - $\beta \in \mathcal{S}_0 \Rightarrow \beta$ *is an unitary Pisot number (Frougny).*

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  - *The first equivalence is not valid in general (Chacon),*
  - $\beta \in \mathcal{S}_0 \Rightarrow \beta$ *is an unitary Pisot number (Frougny).*
- *If* $p > 1$ *and* $\beta \in \mathcal{S}_0$*, then* $\mathbf{u}_\beta$ *and* $0^{-1}\mathbf{u}_\beta$ *are the only infinite LS branches.*

Affine complexity

Theorem

- *The factor complexity of* $\mathbf{u}_\beta$ *is affine* $\Leftrightarrow$ $\mathbf{u}_\beta$ *does not contain any* (a, b)*-maximal factor* $\Leftrightarrow \beta \in \mathcal{S}_0 \Leftrightarrow d_\beta(1) = t_1(0 \cdots 0(t_1 - 1))^\omega$. *Then* $\mathcal{C}(n) = (m + p - 1)n + 1$.
  - *The first equivalence is not valid in general (Chacon),*
  - $\beta \in \mathcal{S}_0 \Rightarrow \beta$ *is an unitary Pisot number (Frougny).*
- *If* $p > 1$ *and* $\beta \in \mathcal{S}_0$, *then* $\mathbf{u}_\beta$ *and* $0^{-1}\mathbf{u}_\beta$ *are the only infinite LS branches.*
- *known result:* $\mathbf{u}_\beta$ *is Sturmian* $\Leftrightarrow p = 1$ *and* $\beta \in \mathcal{S}_0$, *i.e.* $d_\beta(1) = t_1(t_1 - 1)^\omega$.

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THE END