

EXTENSIONS AND RESTRICTIONS OF WYTHOFF'S GAME  
PRESERVING WYTHOFF'S SEQUENCE AS SET OF  $\mathcal{P}$  POSITIONS

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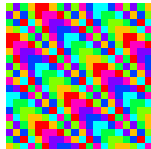
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Michel Rigo (University of Liège)

<http://www.discmath.ulg.ac.be/>

Dynamical Aspects of Numeration Systems, Roma, Feb. 2008



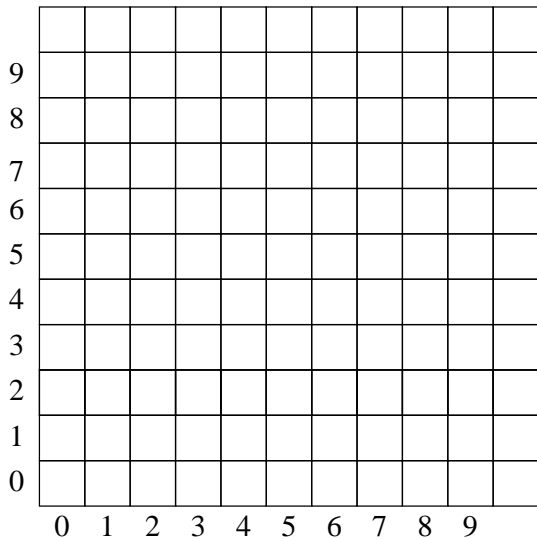
W. A. Wythoff, A modification of the game of Nim,  
*Nieuw Arch. Wisk.* **7** (1907), 199–202.

## RULES OF THE GAME

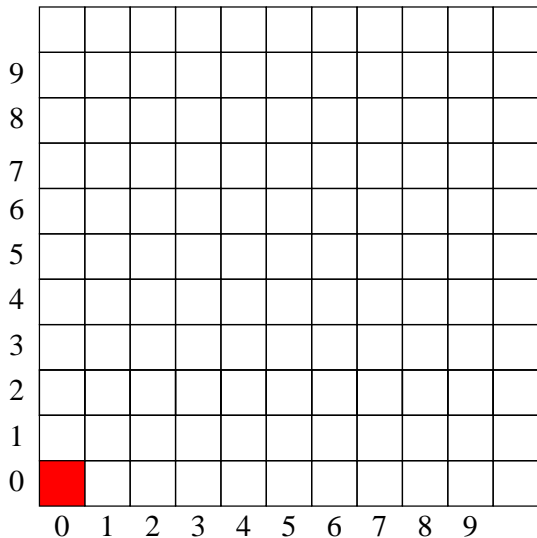
- ▶ Two players play alternatively
- ▶ Two piles of tokens
- ▶ Remove
  - ▶ any positive number of tokens from **one** pile or,
  - ▶ the **same** positive number from the two piles.
- ▶ The one who takes the last token wins the game (**last move wins**).

Set of moves :  $\{(i, 0), i > 0\} \cup \{(0, j), j > 0\} \cup \{(k, k), k > 0\}$

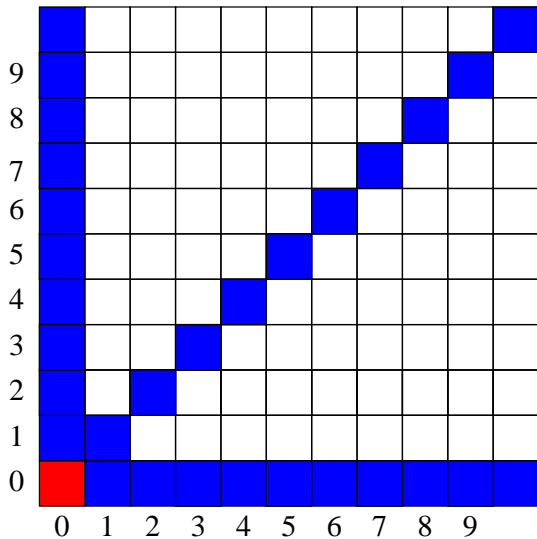
# WYTHOFF'S GAME OR "CATCHING THE QUEEN"



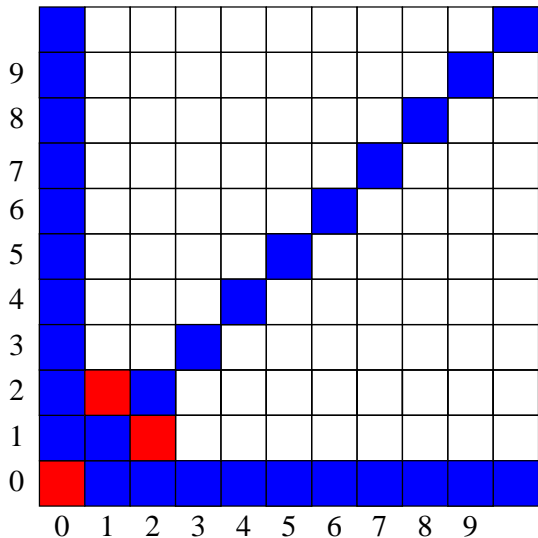
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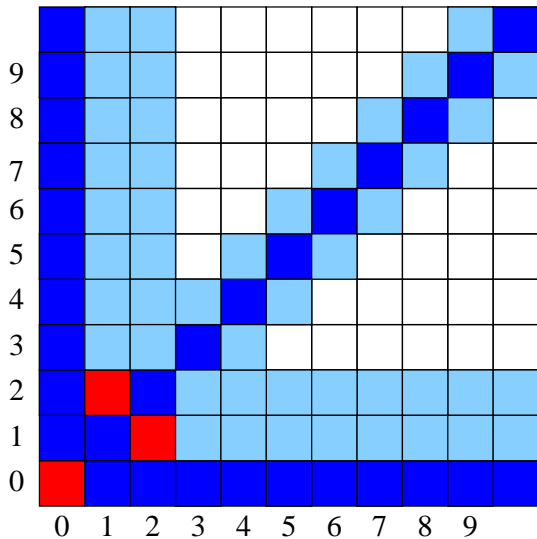
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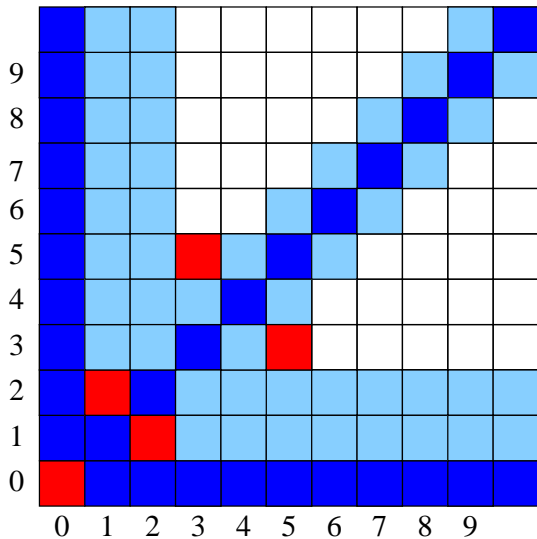
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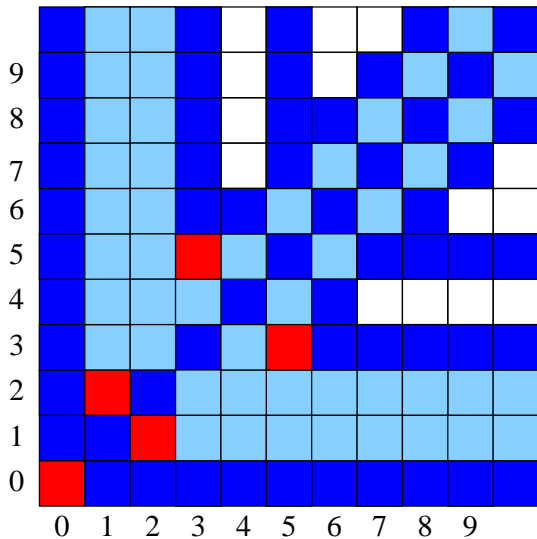


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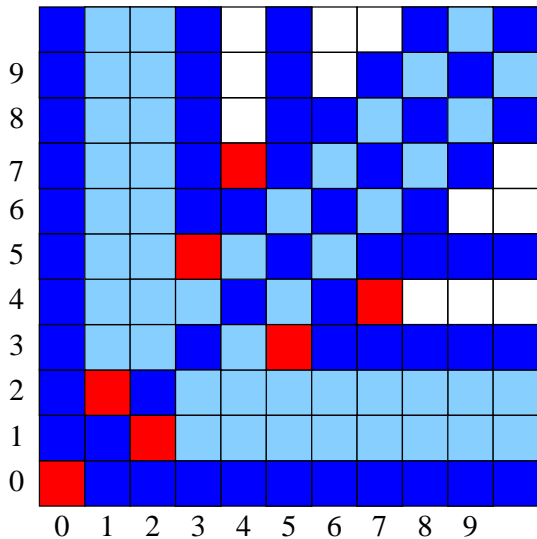




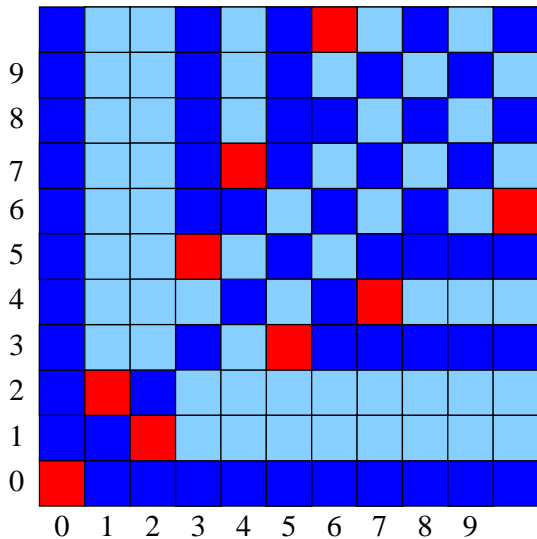
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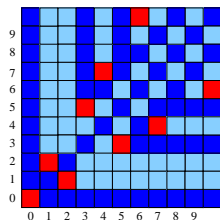
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## WYTHOFF'S GAME OR "CATCHING THE QUEEN"



$(0, 0), (1, 2), (3, 5), (4, 7), (6, 10), \dots$

### P-POSITION

A **P-position** is a position  $q$  from which the *previous* player (moving to  $q$ ) can force a win.

### N-POSITION

A **N-position** is a position  $p$  from which the *actual* player has an option leading ultimately to win the game.

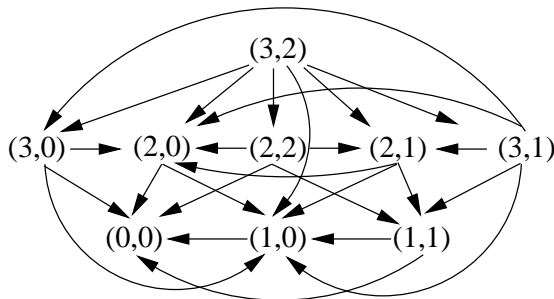
Question : **Are all positions  $\mathcal{N}$  or  $\mathcal{P}$  ?**

# GAME GRAPH

Initial position  $(i_0, j_0)$ , by symmetry, take only  $(i \geq j)$

- ▶ **Vertices** :  $\{(i, j), i \leq i_0, j \leq j_0\}$
- ▶ **Edges** : from each position to all its options :

$$\begin{array}{l} i > 0 \\ j > 0 \\ i, j > 0 \end{array} \left| \begin{array}{l} (i, j) \rightarrow (i-k, j) \\ (i, j) \rightarrow (i, j-k) \\ (i, j) \rightarrow (i-k, j-k) \end{array} \right| \begin{array}{l} k = 1, \dots, i \\ k = 1, \dots, j \\ k = 1, \dots, \min(i, j) \end{array}$$



## REMARK

Due to the rules, the game graph for Wythoff's game is **acyclic**.

## THEOREM [BERGE]

Any finite **acyclic** digraph has a **unique kernel**.

Moreover, this kernel can be obtained efficiently.

## REMINDER/DEFINITION OF A KERNEL

A **kernel** in a graph  $G = (V, E)$  is a subset  $W \subseteq V$

- ▶ **stable** :  $\forall x, y \in W, (x, y) \notin E$
- ▶ **absorbing** :  $\forall x \in V \setminus W, \exists y \in W : (x, y) \in E.$

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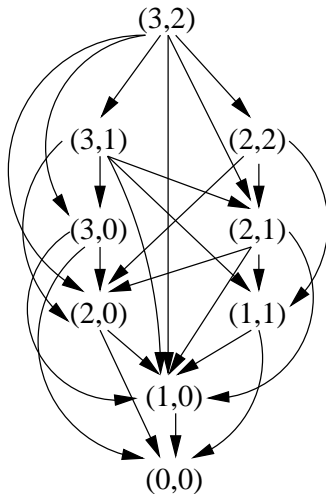
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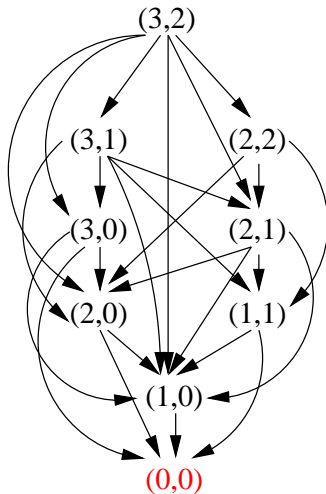
Bottom-Up approach from the sinks  
(they belong to the kernel because it is absorbing)





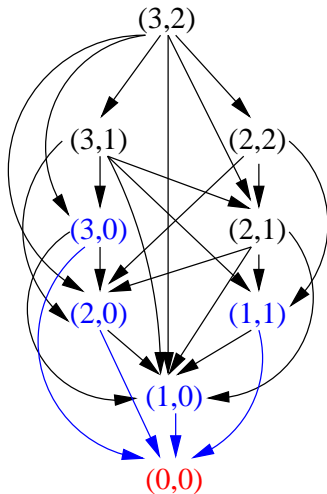
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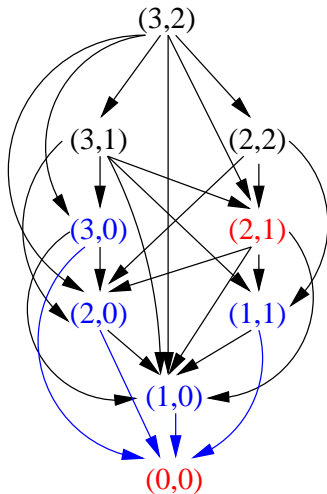
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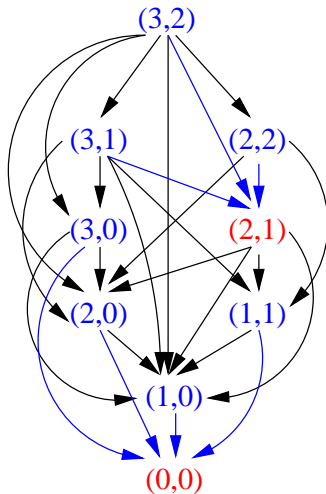
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# GAME GRAPH

For Wythoff's game, its game graph has a unique kernel  $K$ .

- ▶ **stable** : from a position in  $K$ , you always play out of  $K$ ,
- ▶ **absorbing** : from a position outside  $K$ , you can play into  $K$ ,
- ▶  $(0,0)$  has to belong to  $K$ , otherwise  $K$  won't be absorbing.

## COROLLARY

The set of  $\mathcal{P}$ -positions is exactly the kernel  $K$  and all the other positions are  $\mathcal{N}$ -positions.

$$\{\mathcal{P}\text{-positions}\} \supseteq K$$

If  $p$  is a position in  $K$ , then it is a  $\mathcal{P}$ -position because there is a *winning strategy* outside  $K$ .

$$\{\mathcal{P}\text{-positions}\} \subseteq K$$

If  $p$  is a  $\mathcal{P}$ -position not in  $K$ , then there is a move from  $p$  to  $K$ , thus  $p$  is a  $\mathcal{N}$ -position !

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## P-POSITION OF THE WYTHOFF'S GAME I

$$(A_n, B_n)_{n \geq 0} = (0, 0), (1, 2), (3, 5), (4, 7), \dots$$

$$\forall n \geq 0, \quad \begin{cases} A_n = \text{Mex}\{A_i, B_i \mid i < n\} \\ B_n = A_n + n \end{cases}$$

## P-POSITION OF THE WYTHOFF'S GAME II

	1	2	3	4	5	6	7	8	9	10	11	12	13	14	...
$\mathcal{F}$	a	b	a	a	b	a	b	a	a	b	a	a	b	a	

## P-POSITIONS OF THE WYTHOFF'S GAME III

$$(A_n, B_n)_{n \geq 0} = (\lfloor n\tau \rfloor, \lfloor n\tau^2 \rfloor).$$

- ▶ A.S. Fraenkel, How to beat your Wythoff games' opponent on three fronts, *Amer. Math. Monthly* **89** (1982), 353–361.
- ▶ A.S. Fraenkel, Heap games, Numeration systems and Sequences, *Annals of Combinatorics* **2** (1998), 197–210.
- ▶ A.S. Fraenkel, The Raleigh Game, *INTEGERS* (2007).
- ▶ E. Duchêne, M.R., A morphic approach to combinatorial games: the Tribonacci case, to appear in *RAIRO Theoret. Inform. Appl.*
- ▶ E. Duchêne, M.R., A class a cubic Pisot unit games, to appear in *Monat. für Math.*

**Different sets of moves / more piles**



**Different sets of  $\mathcal{P}$ -positions to characterize...**



## OUR GOAL / DUAL QUESTION

Consider extensions or restrictions of Wythoff's game that keep the set of  $\mathcal{P}$ -positions of Wythoff's game **invariant**.

**Characterize the different sets of moves...**

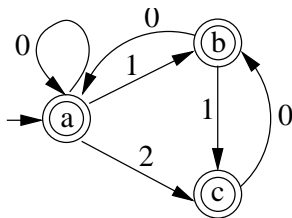


**Same set of  $\mathcal{P}$ -positions as Wythoff's game**

# DURING OUR JOURNEY...

Canonical construction [Cobham'72] : morphisms  $\rightarrow$  automata

$$\varphi : a \mapsto abc, b \mapsto ac, c \mapsto b$$

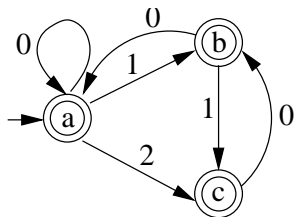


$$\varphi^\omega(a) = abcacbabcbaacabcbacabcbabcacb \dots$$

Consider the language  $L = L(\mathcal{M}) \setminus 0\{0, 1, 2\}^*$ .

**Remark:** Positions in  $\varphi^\omega(a)$  are counted from 1.

Take the words of  $L$  in genealogical order (abstract system)



$n$	$w_n$			$n$	$w_n$		
0	$\epsilon$	a	1	10	200	a	11
1	1	b	2	11	201	c	12
2	2	c	3	12	1000	a	13
3	10	a	4	13	1001	b	14
4	11	c	5	14	1002	c	15
5	20	b	6	15	1010	a	16
6	100	a	7	16	1011	c	17
7	101	b	8	17	1020	b	18
8	102	c	9	18	1100	a	19
9	110	0	10	19	1101	c	20

Not a “positional” system, no sequence behind.

### EXAMPLE :

The 4th letter is a, it corresponds to  $w_3 = 10$ .

Since  $\varphi(a) = abc$ , we consider

$$\begin{cases} w_30 = 100 = w_i \\ w_31 = 101 = w_{i+1} \\ w_32 = 102 = w_{i+2} \end{cases}$$

then the  $(i + 1)$ st,  $(i + 2)$ st,  $(i + 3)$ st letters are a, b, c.

$$\text{rep}_L(i) := w_i, \quad \text{val}_L(w_i) := i$$

## PROPOSITION

Let the  $n$ th letter of  $\varphi^\omega(a)$  be  $\sigma$  and  $w_{n-1}$  be the  $n$ th word in  $L$ . If  $\varphi(\sigma) = x_1 \cdots x_r$ , then  $x_1 \cdots x_r$  appears in  $\varphi^\omega(a)$  in positions  $\text{val}_L(w_{n-1}x_1)+1, \dots, \text{val}_L(w_{n-1}x_r)+1$ .

For Wythoff's game: Fibonacci word  $\mathcal{F}$ ,  $L = 1\{01, 0\}^* \cup \{\varepsilon\}$  and we get the usual Fibonacci system  $\rho_F : \mathbb{N} \rightarrow L$ ,  $\pi_F : L \rightarrow \mathbb{N}$ .

## COROLLARY

- ▶ If the  $n$ th letter in  $\mathcal{F}$  is  $a$  ( $n \geq 1$ ), then this  $a$  produces through  $\varphi$  a factor  $ab$  occupying positions  $\pi_F(\rho_F(n-1)0)+1$  and  $\pi_F(\rho_F(n-1)1)+1$ .
- ▶ If the  $n$ th letter in  $\mathcal{F}$  is  $b$  ( $n \geq 1$ ), then this  $b$  produces through  $\varphi$  a letter  $a$  occupying position  $\pi_F(\rho_F(n-1)0) + 1$ .

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## REMINDER ON FIBONACCI NUMERATION SYSTEM

Fibonacci sequence :  $F_{i+2} = F_{i+1} + F_i$ ,  $F_0 = 1$ ,  $F_1 = 2$

Use **greedy** expansion,  $\dots$ , 21, 13, 8, 5, 3, 2, 1

$n$	$\rho_F(n)$	$n$	$\rho_F(n)$	$n$	$\rho_F(n)$
1	1	8	10000	15	100010
2	10	9	10001	16	100100
3	100	10	10010	17	100101
4	101	11	10100	18	101000
5	1000	12	10101	19	101001
6	1001	13	100000	20	101010
7	1010	14	100001	21	1000000

E. Zeckendorf, Représentation des nombres naturels par une somme des nombres de Fibonacci ou de nombres de Lucas, *Bull. Soc. Roy. Sci. Liège* **41** (1972), 179–182.

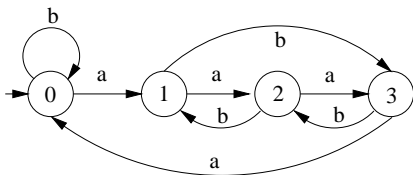
In fact, this is a special case of the following result.

**THEOREM [A. MAES, M.R. '02]**

The set of S-automatic sequences is exactly the set of morphic words.

Take any regular language genealogically ordered  $\oplus$  DFAO

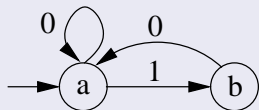
$i$	0	1	2	3	4	5	6	7	8	9	...
$\text{rep}_S(i)$	$\epsilon$	$a$	$b$	$aa$	$ab$	$bb$	$aaa$	$aab$	$abb$	$bbb$	...



01023031200231010123023031203120231002310123010123...

$n$	1	2	3	4	5	6	7	8	9	10	11	12
	$a$	$b$	$a$	$a$	$b$	$a$	$b$	$a$	$a$	$b$	$a$	$a$
$A_i$	1		3	4		6		8	9		11	12
$B_i$		2			5		7			10		
$\rho_F(n-1)$	$\omega$	$\tau$	10	100	101	1000	1001	1010	10000	10001	10010	10100

## P-POSITIONS OF THE WYTHOFF'S GAME IV



For all  $n \geq 1$ , we have

$$A_n = \pi_F(\rho_F(n-1)0) + 1$$

$$B_n = \pi_F(\rho_F(A_n-1)1) + 1.$$



# MORE ?

Can we get a “morphic characterization” of the Wythoff’s matrix ?

$$(P_{i,j})_{i,j \geq 0} = \begin{array}{cccccccccccc} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \dots \\ 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & \\ 0 & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{1} & \\ 0 & 0 & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & \\ \vdots & & & & & & & & & & & \ddots \end{array}$$

Let's try something...

$$\varphi : a \mapsto \begin{array}{|c|c|} \hline a & b \\ \hline c & d \\ \hline \end{array} \quad b \mapsto \begin{array}{|c|} \hline i \\ \hline e \\ \hline \end{array} \quad c \mapsto \boxed{i \mid j} \quad d \mapsto \boxed{i} \quad e \mapsto \boxed{f \mid b}$$

$$f \mapsto \begin{array}{|c|c|} \hline g & b \\ \hline h & d \\ \hline \end{array} \quad g \mapsto \begin{array}{|c|c|} \hline f & b \\ \hline h & d \\ \hline \end{array} \quad h \mapsto \boxed{i \mid m} \quad i \mapsto \begin{array}{|c|c|} \hline i & m \\ \hline h & d \\ \hline \end{array}$$

$$j \mapsto \begin{array}{|c|} \hline k \\ \hline c \\ \hline \end{array} \quad k \mapsto \begin{array}{|c|c|} \hline l & m \\ \hline c & d \\ \hline \end{array} \quad l \mapsto \begin{array}{|c|c|} \hline k & m \\ \hline c & d \\ \hline \end{array} \quad m \mapsto \begin{array}{|c|} \hline i \\ \hline h \\ \hline \end{array}$$

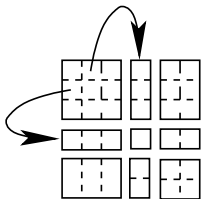
and the coding

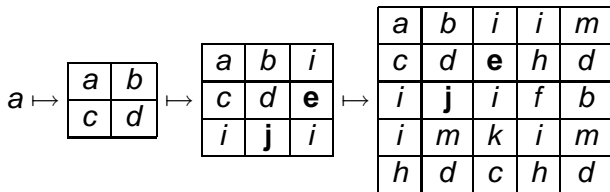
$$\mu : e, g, j, l \mapsto 1, \quad a, b, c, d, f, h, i, k, m \mapsto 0$$

O. Salon, Suites automatiques à multi-indices, *Séminaire de théorie des nombres*, Bordeaux, 1986–1987, exposé 4.

## SHAPE-SYMMETRIC MORPHISM [A. MAES '99]

If  $P$  is the infinite bidimensional picture that is the fixpoint of  $\varphi$ , then for all  $i, j \in \mathbb{N}$ , if  $\varphi(P_{i,j})$  is a block of size  $k \times \ell$  then  $\varphi(P_{j,i})$  is of size  $\ell \times k$

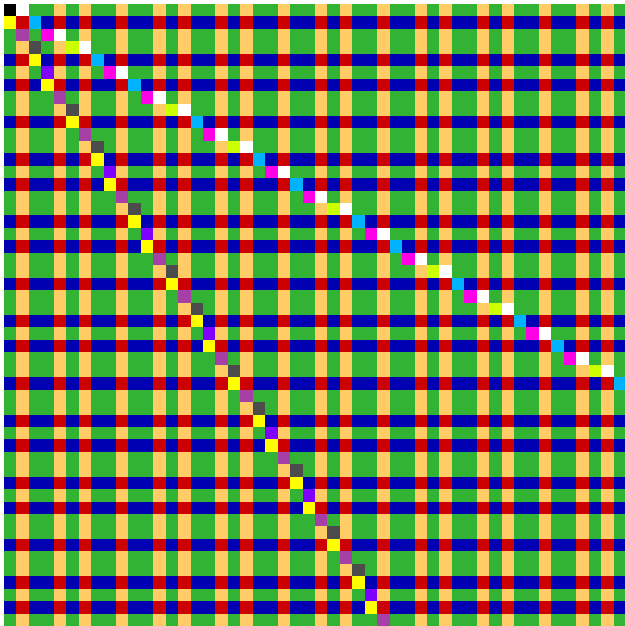




sizes : 1, 2, 3, 5

<i>a</i>	<i>b</i>	<i>i</i>	<i>i</i>	<i>m</i>	<i>i</i>	<i>m</i>	<i>i</i>
<i>c</i>	<i>d</i>	<b><i>e</i></b>	<i>h</i>	<i>d</i>	<i>h</i>	<i>d</i>	<i>h</i>
<i>i</i>	<b><i>j</i></b>	<i>i</i>	<i>f</i>	<i>b</i>	<i>i</i>	<i>m</i>	<i>i</i>
<i>i</i>	<i>m</i>	<i>k</i>	<i>i</i>	<i>m</i>	<b><i>g</i></b>	<i>b</i>	<i>i</i>
<i>h</i>	<i>d</i>	<i>c</i>	<i>h</i>	<i>d</i>	<i>h</i>	<i>d</i>	<b><i>e</i></b>
<i>i</i>	<i>m</i>	<i>i</i>	<b><i>l</i></b>	<i>m</i>	<i>i</i>	<i>m</i>	<i>i</i>
<i>h</i>	<i>d</i>	<i>h</i>	<i>c</i>	<i>d</i>	<i>h</i>	<i>d</i>	<i>h</i>
<i>i</i>	<i>m</i>	<i>i</i>	<i>i</i>	<b><i>j</i></b>	<i>i</i>	<i>m</i>	<i>i</i>

size : 8,...



# MORPHISMS $\rightarrow$ AUTOMATA

We can do the same as for the unidimensional case :  
Automaton with input alphabet

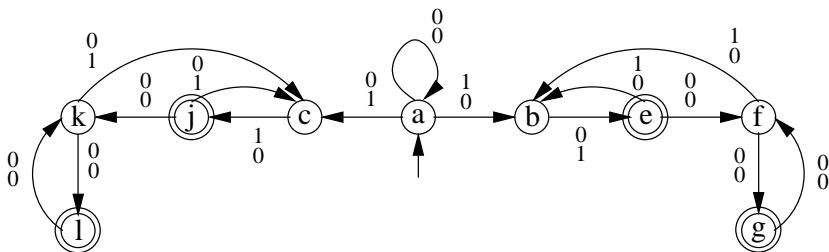
$$\left\{ \begin{pmatrix} 0 \\ 0 \end{pmatrix}, \begin{pmatrix} 1 \\ 0 \end{pmatrix}, \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \begin{pmatrix} 1 \\ 1 \end{pmatrix} \right\}$$

$$\varphi(r) = \begin{array}{|c|c|} \hline s & t \\ \hline u & v \\ \hline \end{array}, \quad \begin{array}{|c|c|} \hline s & t \\ \hline \end{array}, \quad \begin{array}{|c|} \hline s \\ \hline u \\ \hline \end{array} \quad \text{or} \quad \begin{array}{|c|} \hline s \\ \hline \end{array}$$

we have transitions like

$$r \xrightarrow{\begin{pmatrix} 0 \\ 0 \end{pmatrix}} s, \quad r \xrightarrow{\begin{pmatrix} 1 \\ 0 \end{pmatrix}} t, \quad r \xrightarrow{\begin{pmatrix} 0 \\ 1 \end{pmatrix}} u, \quad r \xrightarrow{\begin{pmatrix} 1 \\ 1 \end{pmatrix}} v.$$

We get (after trimming useless part with four states)



This automaton accepts the words

$$\begin{pmatrix} 0w_1 \cdots w_\ell \\ w_1 \cdots w_\ell 0 \end{pmatrix} \text{ and } \begin{pmatrix} w_1 \cdots w_\ell 0 \\ 0w_1 \cdots w_\ell \end{pmatrix}$$

where  $w_1 \cdots w_\ell$  is a valid  $F$ -representation **ending with an even number of zeroes.**



Such a characterization is well-known, but differs from the one we get previously...

## REMINDER

For all  $n \geq 1$ , we have

$$\begin{aligned}A_n &= \pi_F(\rho_F(n-1)0) + 1 \\B_n &= \pi_F(\rho_F(A_n-1)1) + 1.\end{aligned}$$

It is hopefully the same, but **why** ?

- First case :  $\rho_F(n-1) = u0$

$$\rho_F(A_n) = \rho_F(\underbrace{\pi_F(\rho_F(n-1)0)}_{u00} + 1) = u01 \text{ no zero}$$

$$\rho_F(A_n - 1) = u00 \text{ and}$$

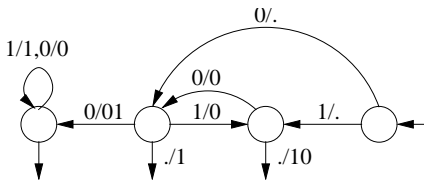
$$\rho_F(B_n) = \rho_F(\underbrace{\pi_F(\rho_F(A_n - 1)1)}_{u001} + 1) = u010 \text{ one zero}$$

- Second case :  $\rho_F(n-1) = u01$

$$\rho_F(A_n) = \rho_F(\underbrace{\pi_F(\rho_F(n-1)0)}_{u010} + 1) = "u011" \dots$$

Normalize  $u011$  or look for the successor of  $u010$

Use the transducer (R to L) computing the successor  
 [Frougny'97]



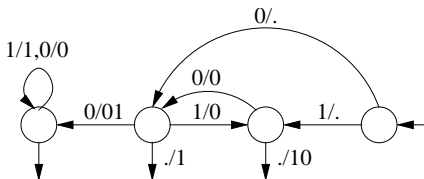
$10 \rightarrow 100$ , 2 zeroes

$\underbrace{x10(01)^n 010}_u \rightarrow x101(00)^n 00$   $2n + 2$  zeroes,  $n \geq 0$

$\underbrace{1(01)^n 010}_u \rightarrow 100(00)^n 00$   $2n + 4$  zeroes,  $n \geq 0$

$$\rho_F(A_n - 1) = u010 \text{ and}$$

$$\rho_F(B_n) = \rho_F(\underbrace{\pi_F(\rho_F(A_n - 1)1)}_{u0101} + 1) = "u0102" \dots$$



$$101 \rightarrow 1000, \quad 3 \text{ zeroes}$$

$$\underbrace{x10(01)^n}_{u} 0101 \rightarrow x101(00)^n 000 \quad 2n + 3 \text{ zeroes}, \quad n \geq 0$$

$$\underbrace{1(01)^n}_{u} 0101 \rightarrow 100(00)^n 000 \quad 2n + 5 \text{ zeroes}, \quad n \geq 0$$

Conclusion : "A<sub>n</sub> even number of zeroes, B<sub>n</sub> one more", OK

# EXTENSION PRESERVING SET OF $\mathcal{P}$ -POSITIONS

To decide whether or not a move can be adjoined to Wythoff's game without changing the set  $K$  of  $\mathcal{P}$ -positions, it suffices to check that it does not change the stability property  $K$ .

Remark : absorbing property holds true whatever the adjoined move is.

## CONSEQUENCE

A move  $(i, j)$  can be added IFF it prevents to move from a  $\mathcal{P}$ -position to another  $\mathcal{P}$ -position.

In other words, a necessary and sufficient condition for a move  $(i, j)_{i < j}$  to be adjoined is that it does not belong to

$$\{(A_n - A_m, B_n - B_m) : n > m \geq 0\} \cup \{(A_n - B_m, B_n - A_m) : n > m \geq 0\}$$

Thanks to the previous characterizations of  $A_n$ ,  $B_m$ ,

## PROPOSITION

A move  $(i, j)_{i < j}$  can be adjoined to without changing the set of  $\mathcal{P}$ -positions IFF

$$(i, j) \neq (\lfloor n\tau \rfloor - \lfloor m\tau \rfloor, \lfloor n\tau^2 \rfloor - \lfloor m\tau^2 \rfloor) \quad \forall n > m \geq 0$$

and

$$(i, j) \neq (\lfloor n\tau \rfloor - \lfloor m\tau^2 \rfloor, \lfloor n\tau^2 \rfloor - \lfloor m\tau \rfloor) \quad \forall n > m \geq 0$$

For all  $i, j \geq 0$ ,  $W_{i,j} = 1$  IFF Wythoff's game with the adjoined move  $(i, j)$  has Wythoff's sequence as set of  $\mathcal{P}$ -positions,

$$(W_{i,j})_{i,j \geq 0} = \begin{array}{cccccccccccc} 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & 0 & \dots \\ 0 & 0 & \mathbf{1} & 0 & \mathbf{1} & 0 & 0 & \mathbf{1} & 0 & \mathbf{1} & 0 & \\ 0 & \mathbf{1} & 0 & \mathbf{1} & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & \\ 0 & 0 & \mathbf{1} & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & \mathbf{1} & \\ 0 & \mathbf{1} & 0 & 0 & 0 & 0 & \mathbf{1} & \mathbf{1} & 0 & \mathbf{1} & 0 & \\ 0 & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & \mathbf{1} & 0 & 0 & \\ 0 & 0 & \mathbf{1} & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & \mathbf{1} & \\ 0 & \mathbf{1} & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & \\ 0 & \mathbf{1} & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & 0 & 0 & \\ 0 & 0 & 0 & \mathbf{1} & 0 & 0 & \mathbf{1} & 0 & 0 & 0 & 0 & \\ \vdots & & & & & & & & & & & \ddots \end{array}$$

## COROLLARY

Let  $I \subseteq \mathbb{N}$ . Wythoff's game with adjoined moves

$$\{(x_i, y_i) : i \in I, x_i, y_i \in \mathbb{N}\}$$

has the same sequence  $(A_n, B_n)$  as set of  $\mathcal{P}$ -positions

IFF

$W_{x_i, y_i} \neq 1$  for all  $i \in I$ .



# ARE WE DONE ? Complexity issue

We investigate **tractable extensions** of Wythoff's game, we also need to **test these conditions in polynomial time**. And the winner can consummate a win in at most an exponential number of moves.

## MANY "EFFORTS" LEAD TO THIS

For any pair  $(i, j)$  of positive integers, we have  $W_{i,j} = 1$  if and only if one the three following properties is satisfied :

- ▶  $(\rho_F(i-1), \rho_F(j-1)) = (u0, u01)$  for any valid  $F$ -representation  $u$  in  $\{0, 1\}^*$ .
- ▶  $(\rho_F(i-2), \rho_F(j-2)) = (u0, u01)$  for any valid  $F$ -representation  $u$  in  $\{0, 1\}^*$ .
- ▶  $(\rho_F(j - A_i - 2), \rho_F(j - A_i - 2 + i)) = (u1, u'0)$  for any two valid  $F$ -representations  $u$  and  $u'$  in  $\{0, 1\}^*$ .

# MORPHIC CHARACTERIZATION OF $W...$ IN PROGRESS

$$\psi : a \mapsto \begin{array}{|c|c|} \hline a & b \\ \hline c & d \\ \hline \end{array} \quad b \mapsto \begin{array}{|c|} \hline e \\ \hline f \\ \hline \end{array} \quad c \mapsto \begin{array}{|c|c|} \hline e & h \\ \hline \end{array} \quad d \mapsto \begin{array}{|c|} \hline i \\ \hline \end{array} \quad e \mapsto \begin{array}{|c|c|} \hline j & k \\ \hline l & m \\ \hline \end{array}$$

$$f \mapsto \begin{array}{|c|c|} \hline g & b \\ \hline \end{array} \quad g \mapsto \begin{array}{|c|c|} \hline y & b \\ \hline o & t \\ \hline \end{array} \quad h \mapsto \begin{array}{|c|} \hline z \\ \hline c \\ \hline \end{array} \quad i \mapsto \begin{array}{|c|c|} \hline i & n \\ \hline o & d \\ \hline \end{array}$$

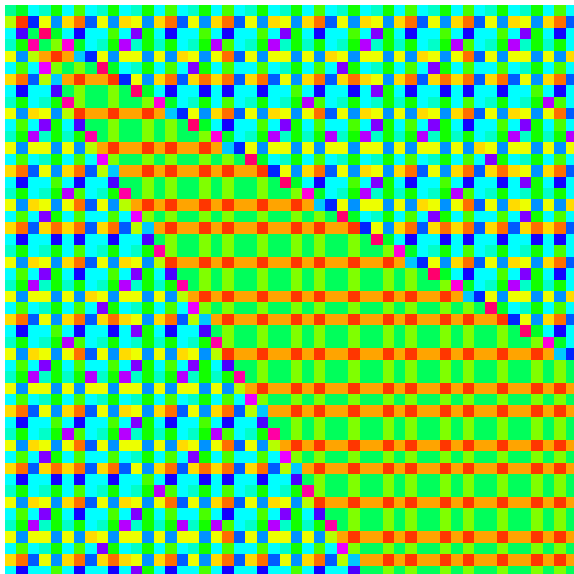
$$j \mapsto \begin{array}{|c|c|} \hline e & p \\ \hline q & r \\ \hline \end{array} \quad k \mapsto \begin{array}{|c|} \hline e \\ \hline s \\ \hline \end{array} \quad l \mapsto \begin{array}{|c|c|} \hline e & u \\ \hline \end{array} \quad m \mapsto \begin{array}{|c|} \hline e \\ \hline \end{array}$$

$$n \mapsto \begin{array}{|c|} \hline i \\ \hline o \\ \hline \end{array} \quad o \mapsto \begin{array}{|c|c|} \hline i & n \\ \hline \end{array} \quad p \mapsto \begin{array}{|c|} \hline e \\ \hline q \\ \hline \end{array} \quad q \mapsto \begin{array}{|c|c|} \hline e & p \\ \hline \end{array} \quad r \mapsto \begin{array}{|c|} \hline e \\ \hline \end{array}$$

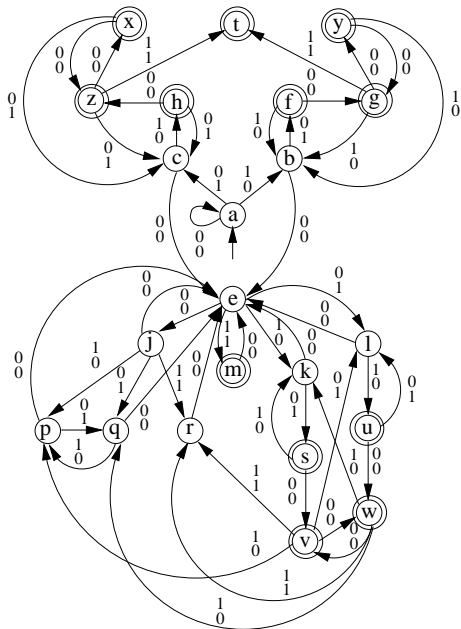
$$s \mapsto \begin{array}{|c|c|} \hline v & k \\ \hline \end{array} \quad t \mapsto \begin{array}{|c|} \hline i \\ \hline \end{array} \quad u \mapsto \begin{array}{|c|} \hline w \\ \hline l \\ \hline \end{array} \quad v \mapsto \begin{array}{|c|c|} \hline w & p \\ \hline l & r \\ \hline \end{array} \quad w \mapsto \begin{array}{|c|c|} \hline v & k \\ \hline q & r \\ \hline \end{array}$$

$$x \mapsto \begin{array}{|c|c|} \hline z & n \\ \hline c & d \\ \hline \end{array} \quad y \mapsto \begin{array}{|c|c|} \hline g & b \\ \hline o & d \\ \hline \end{array} \quad z \mapsto \begin{array}{|c|c|} \hline x & n \\ \hline c & t \\ \hline \end{array}$$

and the coding  $\nu : a, b, c, d, e, i, j, k, l, n, o, p, q, r \mapsto 0$   
 $f, g, h, m, s, t, u, v, w, x, y, z \mapsto 1$ .



# Corresponding automaton



# SOME OF THE MACHINERY BEHIND



## LEMMA

Let  $\mathcal{F}_n$  be the prefix of  $\mathcal{F}$  of length  $n$ .

For any finite factor  $bu\bar{a}$  occurring in  $\mathcal{F}$  with  $|u| = n$ ,  
we have  $|u|_a = |\mathcal{F}_n|_a$  and  $|u|_b = |\mathcal{F}_n|_b$ .

## EXAMPLE

Take  $u = aabaab$ ,  $bu\bar{a}$  of length 8 starts in  $\mathcal{F}$  from position 7.

$\mathcal{F}_6 = abaaba$  is a permutation of  $u$ .

$$\mathcal{F} = \underbrace{abaaba}_{\mathcal{F}_6} b \overbrace{aabaab}^{bu\bar{a}} a \text{baababaaba} \dots$$

$\underbrace{\hspace{1.5cm}}_u$

Proof : algebraic

## LEMMA

Let  $n \geq 1$  be such that  $B_{n+1} - B_n = 2$ . Then  $\rho_F(B_n - 1)$  ends with 101.

Proof : Morphic structure of  $\mathcal{F}$

## PROPOSITION

$$\{(A_j - A_i, B_j - B_i) \mid j > i \geq 0\} = \{(A_n, B_n) \mid n > 0\} \\ \cup \{(A_n + 1, B_n + 1) \mid n > 0\}$$

Proof : Density of the  $\{n\tau\}$ 's in  $[0, 1]$

## LEMMA

Let  $u1 \in \{0, 1\}^*$  be a valid  $F$ -representation. If  $\rho_F(\pi_F(u1) + n)1$  is also a valid  $F$ -representation, then

$$\pi_F(\rho_F(\pi_F(u1) + n)1) = \pi_F(u00) + \pi_F(\rho_F(n - 1)0) + 4.$$

Otherwise,  $\rho_F(\pi_F(u1) + n)1$  is not a valid  $F$ -representation and

$$\pi_F(\rho_F(\pi_F(u1) + n)0) = \pi_F(u00) + \pi_F(\rho_F(n)0) + 2.$$

Proof : Morphic structure of  $\mathcal{F}$

## THEOREM

Let  $i, j$  be such that  $A_j - B_i = n > 0$ . We have

$$B_j - A_i = B_i + A_n + 1.$$



## THEOREM

There is no redundant move in Wythoff's game. In particular, if any move is removed, then the set of  $\mathcal{P}$ -positions changes.